

An Algorithm for Solving Non-Symmetric Saddle-Point Systems

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OUTLINE

Motivation: Fictitious domain method

Algorithm PSCM: Schur complement method + Null-space method

Inner solver: Projected BiCGSTAB

Preconditioning: Hierarchical multigrid

Singular matrices: Poisson-like solver based on circulants

Numerical experiments

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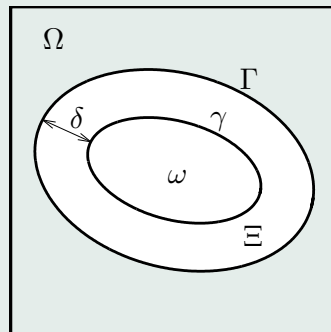
MODEL PROBLEM 1: Dirichlet problem

$$-\Delta u = f \quad \text{on } \omega \quad (1)$$

$$u = g \quad \text{in } \gamma \equiv \partial\omega \quad (2)$$

Fictitious domain method (FDM):

PDE (1) is solved on the fictitious domain $\Omega, \bar{\omega} \subset \Omega$, with a simple geometry. The corresponding stiffness matrix \mathbf{A} is structured. The original boundary conditions (2) on γ are enforced by Lagrange multipliers or control variables.



Classical FDM with $\Gamma = \gamma$

$$\begin{pmatrix} \mathbf{A} & \mathbf{B}_\gamma^\top \\ \mathbf{B}_\gamma & \mathbf{0} \end{pmatrix} \begin{pmatrix} \mathbf{u} \\ \boldsymbol{\lambda}_\gamma \end{pmatrix} = \begin{pmatrix} \mathbf{f} \\ \mathbf{g} \end{pmatrix}$$

Smooth FDM with $\Gamma \neq \gamma$

$$\begin{pmatrix} \mathbf{A} & \mathbf{B}_\Gamma^\top \\ \mathbf{B}_\gamma & \mathbf{0} \end{pmatrix} \begin{pmatrix} \mathbf{u} \\ \boldsymbol{\lambda}_\Gamma \end{pmatrix} = \begin{pmatrix} \mathbf{f} \\ \mathbf{g} \end{pmatrix}$$

MODEL PROBLEM 2: Signorini problem

$$-\Delta u = f \quad \text{on } \omega \quad (3)$$

$$u - g \geq 0, \quad \frac{\partial u}{\partial n_\gamma} \geq 0, \quad (u - g) \frac{\partial u}{\partial n_\gamma} = 0 \quad \text{in } \gamma \equiv \partial\omega \quad (4)$$

FDM formulation uses the non-differentiable max-function to express BC (4):

$$\left. \begin{aligned} \mathbf{A}\mathbf{u} + \mathbf{B}_\Gamma \boldsymbol{\lambda}_\Gamma &= \mathbf{f} \\ \mathbf{C}_{\gamma,i} \mathbf{u} &= \max \{0, \mathbf{C}_{\gamma,i} \mathbf{u} - \rho(\mathbf{B}_{\gamma,i} \mathbf{u} - \mathbf{g}_i)\}, \quad i = 1, \dots, m \end{aligned} \right\} \quad (5)$$

where $\mathbf{B}_{\gamma,i}$, $\mathbf{B}_{\Gamma,i}$ and $\mathbf{C}_{\gamma,i}$ are rows of Dirichlet and Neumann trace matrices, respectively.

The equations (5) can be solved by the semi-smooth Newton method, in which

$$\mathbf{Jacobian} = \begin{pmatrix} \mathbf{A} & \mathbf{B}_\Gamma^\top \\ \partial G(\mathbf{u}) & \mathbf{0} \end{pmatrix}$$

is determined by the generalized derivative $\partial G(\mathbf{u})$.

MODEL PROBLEM 2: Newton method = Active set algorithm

(0) Set $k := 1$, $\rho > 0$, $\varepsilon_u > 0$, $\mathbf{u}^{(0)} \in \mathbb{R}^n$, $\boldsymbol{\lambda}^{(0)} \in \mathbb{R}^m$.

(1) Define the inactive and active sets by:

$$\begin{aligned}\mathcal{I}^k &:= \{i : \mathbf{C}_{\gamma,i} \mathbf{u}^{k-1} - \rho(\mathbf{B}_{\gamma,i} \mathbf{u}^{k-1} - \mathbf{g}_i) \leq 0\} \\ \mathcal{A}^k &:= \{i : \mathbf{C}_{\gamma,i} \mathbf{u}^{k-1} - \rho(\mathbf{B}_{\gamma,i} \mathbf{u}^{k-1} - \mathbf{g}_i) > 0\}\end{aligned}$$

(2) Solve:

$$\begin{pmatrix} \mathbf{A} & \mathbf{B}_{\Gamma}^{\top} \\ \mathbf{B}_{\gamma, \mathcal{A}^k} & \mathbf{0} \\ \mathbf{C}_{\gamma, \mathcal{I}^k} & \mathbf{0} \end{pmatrix} \begin{pmatrix} \mathbf{u}^k \\ \boldsymbol{\lambda}_{\Gamma}^k \end{pmatrix} = \begin{pmatrix} \mathbf{f} \\ \mathbf{g}_{\mathcal{A}^k} \\ \mathbf{0} \end{pmatrix}$$

(3) If $\|\mathbf{u}^k - \mathbf{u}^{k-1}\| / \|\mathbf{u}^k\| \leq \varepsilon_u$, return $\mathbf{u} := \mathbf{u}^k$.

(4) Set $k := k + 1$, and go to step (1).

Remark: The mixed Dirichlet-Neumann problem is solved in each Newton step, that is described by the non-symmetric saddle-point system.

FORMULATION: Non-symmetric saddle-point system

$$\begin{pmatrix} \mathbf{A} & \mathbf{B}_1^\top \\ \mathbf{B}_2 & \mathbf{0} \end{pmatrix} \begin{pmatrix} \mathbf{u} \\ \lambda \end{pmatrix} = \begin{pmatrix} \mathbf{f} \\ \mathbf{g} \end{pmatrix}$$

General assumptions

\mathbf{A} ... **non-symmetric** ($n \times n$)-matrix
... **singular** with $p = \dim \text{Ker } \mathbf{A}$

$\mathbf{B}_1, \mathbf{B}_2$... **full rank** ($m \times n$)-matrices
... **$\mathbf{B}_1 \neq \mathbf{B}_2$**

Special FDM assumptions

- n is large ($n = 4198401$)
- $m \ll n$ ($m = 360$)
- $p \ll m$ ($p = 1$)
- \mathbf{A} is **structured** so that actions of \mathbf{A}^\dagger or (\mathbf{A}^{-1}) are "cheap"
- $\mathbf{B}_1, \mathbf{B}_2$ are highly **sparse** so that their actions are "cheap"

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ALGORITHMS based on the Schur complement reduction

Case 1: **A** non-singular, symmetric

Case 2: **A** non-singular, non-symmetric

Case 3: **A** singular, symmetric

Case 4: **A** singular, non-symmetric

Case 1: \mathbf{A} non-singular, symmetric

$$\begin{aligned} \begin{pmatrix} \mathbf{A} & \mathbf{B}^\top \\ \mathbf{B} & \mathbf{0} \end{pmatrix} \begin{pmatrix} \mathbf{u} \\ \boldsymbol{\lambda} \end{pmatrix} &= \begin{pmatrix} \mathbf{f} \\ \mathbf{g} \end{pmatrix} &\implies \mathbf{u} &= \mathbf{A}^{-1}(\mathbf{f} - \mathbf{B}^\top \boldsymbol{\lambda}) \\ & &\implies \underbrace{\mathbf{B}\mathbf{A}^{-1}\mathbf{B}^\top}_{\text{negative Schur complement } \mathbf{S}} \boldsymbol{\lambda} &= \mathbf{B}\mathbf{A}^{-1}\mathbf{f} - \mathbf{g} \end{aligned}$$

Algorithm

- 1° Assemble $\mathbf{d} := \mathbf{B}\mathbf{A}^{-1}\mathbf{f} - \mathbf{g}$.
- 2° Solve iteratively $\mathbf{S}\boldsymbol{\lambda} = \mathbf{d}$ with $\mathbf{S} := \mathbf{B}\mathbf{A}^{-1}\mathbf{B}^\top$.
- 3° Assemble $\mathbf{u} := \mathbf{A}^{-1}(\mathbf{f} - \mathbf{B}^\top \boldsymbol{\lambda})$.

If \mathbf{A} is positive defined, then CGM can be used.

Matrix-vector products $\mathbf{S}\boldsymbol{\mu}$ are performed by:

$$\mathbf{S}\boldsymbol{\mu} := \left(\mathbf{B} \left(\mathbf{A}^{-1} \left(\mathbf{B}^\top \boldsymbol{\mu} \right) \right) \right)$$

Case 2: \mathbf{A} non-singular, non-symmetric

$$\begin{aligned} \begin{pmatrix} \mathbf{A} & \mathbf{B}_1^\top \\ \mathbf{B}_2 & \mathbf{0} \end{pmatrix} \begin{pmatrix} \mathbf{u} \\ \lambda \end{pmatrix} &= \begin{pmatrix} \mathbf{f} \\ \mathbf{g} \end{pmatrix} &\implies \mathbf{u} &= \mathbf{A}^{-1}(\mathbf{f} - \mathbf{B}_1^\top \lambda) \\ & &\implies \underbrace{\mathbf{B}_2 \mathbf{A}^{-1} \mathbf{B}_1^\top}_{\text{negative Schur complement } \mathbf{S}} \lambda &= \mathbf{B}_2 \mathbf{A}^{-1} \mathbf{f} - \mathbf{g} \end{aligned}$$

Algorithm is analogous.

- an iterative method for non-symmetric matrices is required (GMRES, BiCG, BiCGSTAB, ...)

$$\mathcal{A} := \begin{pmatrix} \mathbf{A} & \mathbf{B}_1^\top \\ \mathbf{B}_2 & \mathbf{0} \end{pmatrix} = \begin{pmatrix} \mathbf{I} & \mathbf{0} \\ \mathbf{B}_2 \mathbf{A}^{-1} & \mathbf{I} \end{pmatrix} \begin{pmatrix} \mathbf{A} & \mathbf{B}_1^\top \\ \mathbf{0} & -\mathbf{S} \end{pmatrix}$$

Theorem 1 Let \mathbf{A} be non-singular. Then \mathcal{A} is invertible **iff** \mathbf{S} is invertible.

Case 3: \mathbf{A} singular, symmetric

- a generalized inverse \mathbf{A}^\dagger satisfying $\mathbf{A} = \mathbf{A}\mathbf{A}^\dagger\mathbf{A}$
- an $(n \times p)$ -matrix \mathbf{N} whose columns span $\text{Ker } \mathbf{A}$

$$\mathbf{A}\mathbf{u} + \mathbf{B}^\top \boldsymbol{\lambda} = \mathbf{f} \quad \iff \quad \mathbf{f} - \mathbf{B}^\top \boldsymbol{\lambda} \in \text{Im } \mathbf{A} \perp \text{Ker } \mathbf{A}$$

$$\Downarrow$$

$$\mathbf{u} = \mathbf{A}^\dagger(\mathbf{f} - \mathbf{B}^\top \boldsymbol{\lambda}) + \mathbf{N}\boldsymbol{\alpha}$$

$$\Downarrow$$

$$\mathbf{N}^\top(\mathbf{f} - \mathbf{B}^\top \boldsymbol{\lambda}) = \mathbf{0}$$

&

$$\mathbf{B}\mathbf{u} = \mathbf{g}$$

The reduced system:

$$\begin{pmatrix} \mathbf{B}\mathbf{A}^\dagger\mathbf{B}^\top & -\mathbf{B}\mathbf{N} \\ -\mathbf{N}^\top\mathbf{B}^\top & \mathbf{0} \end{pmatrix} \begin{pmatrix} \boldsymbol{\lambda} \\ \boldsymbol{\alpha} \end{pmatrix} = \begin{pmatrix} \mathbf{B}\mathbf{A}^\dagger\mathbf{f} - \mathbf{g} \\ -\mathbf{N}^\top\mathbf{f} \end{pmatrix}$$

$$\Downarrow$$

$$\mathbf{B}\mathbf{A}^\dagger\mathbf{B}^\top \boldsymbol{\lambda} - \mathbf{B}\mathbf{N}\boldsymbol{\alpha} = \mathbf{B}\mathbf{A}^\dagger\mathbf{f} - \mathbf{g}$$

If \mathbf{A} is positive semidefinite, then it corresponds to the algebra in FETI DDM.

Case 4: \mathbf{A} singular, non-symmetric

- a generalized inverse \mathbf{A}^\dagger
- columns of $(n \times p)$ -matrices \mathbf{N} , \mathbf{M} span $\text{Ker } \mathbf{A}$, $\text{Ker } \mathbf{A}^\top$, respectively

$$\mathbf{A}\mathbf{u} + \mathbf{B}_1^\top \boldsymbol{\lambda} = \mathbf{f} \quad \iff \quad \mathbf{f} - \mathbf{B}_1^\top \boldsymbol{\lambda} \in \text{Im } \mathbf{A} \perp \text{Ker } \mathbf{A}^\top$$

$$\Downarrow$$

$$\mathbf{u} = \mathbf{A}^\dagger(\mathbf{f} - \mathbf{B}_1^\top \boldsymbol{\lambda}) + \mathbf{N}\boldsymbol{\alpha}$$

$$\Downarrow$$

$$\mathbf{M}^\top(\mathbf{f} - \mathbf{B}_1^\top \boldsymbol{\lambda}) = \mathbf{0}$$

&

The reduced system:

$$\mathbf{B}_2\mathbf{u} = \mathbf{g}$$

$$\begin{pmatrix} \mathbf{B}_2\mathbf{A}^\dagger\mathbf{B}_1^\top & -\mathbf{B}_2\mathbf{N} \\ -\mathbf{M}^\top\mathbf{B}_1^\top & \mathbf{0} \end{pmatrix} \begin{pmatrix} \boldsymbol{\lambda} \\ \boldsymbol{\alpha} \end{pmatrix} = \begin{pmatrix} \mathbf{B}_2\mathbf{A}^\dagger\mathbf{f} - \mathbf{g} \\ -\mathbf{M}^\top\mathbf{f} \end{pmatrix}$$

$$\Downarrow$$

$$\mathbf{B}_2\mathbf{A}^\dagger\mathbf{B}_1^\top \boldsymbol{\lambda} - \mathbf{B}_2\mathbf{N}\boldsymbol{\alpha} = \mathbf{B}_2\mathbf{A}^\dagger\mathbf{f} - \mathbf{g}$$

Theorem 2 The saddle-point matrix $\mathcal{A} := \begin{pmatrix} \mathbf{A} & \mathbf{B}_1^\top \\ \mathbf{B}_2 & \mathbf{0} \end{pmatrix}$ is invertible **iff**

$$\left. \begin{array}{l} \mathbf{B}_1 \text{ has full row-rank} \\ \text{Ker } \mathbf{A} \cap \text{Ker } \mathbf{B}_2 = \{\mathbf{0}\} \\ \mathbf{A} \text{ Ker } \mathbf{B}_2 \cap \text{Im } \mathbf{B}_1^\top = \{\mathbf{0}\} \end{array} \right\} \text{ (NSC)}$$

Remark: The third equality is equivalent to the MinMax condition that is well-known in the continuous setting:

$$\exists C > 0 : \min_{\mathbf{u} \in \text{Ker } \mathbf{B}_2, \mathbf{u} \neq \mathbf{0}} \max_{\mathbf{v} \in \text{Ker } \mathbf{B}_1, \mathbf{v} \neq \mathbf{0}} \frac{\mathbf{v}^\top \mathbf{A} \mathbf{u}}{\|\mathbf{v}\| \|\mathbf{u}\|} \geq C$$

The generalized Schur complement: the matrix of the reduced system

$$\mathcal{S} := \begin{pmatrix} -\mathbf{B}_2 \mathbf{A}^\dagger \mathbf{B}_1^\top & \mathbf{B}_2 \mathbf{N} \\ \mathbf{M}^\top \mathbf{B}_1^\top & \mathbf{0} \end{pmatrix}$$

Theorem 3 The following three statements are equivalent:

- The necessary and sufficient condition (NSC) holds.
- \mathcal{A} is invertible.
- \mathcal{S} is invertible.

Remark: The generalized Schur complement \mathcal{S} is not defined uniquely.

Second step of the algorithm = Null-space method:

$$\begin{pmatrix} \mathbf{F} & \mathbf{G}_1^\top \\ \mathbf{G}_2 & \mathbf{0} \end{pmatrix} \begin{pmatrix} \boldsymbol{\lambda} \\ \boldsymbol{\alpha} \end{pmatrix} = \begin{pmatrix} \mathbf{d} \\ \mathbf{e} \end{pmatrix}$$

Two orthogonal projectors \mathbf{P}_1 and \mathbf{P}_2 onto $\text{Ker } \mathbf{G}_1$ and $\text{Ker } \mathbf{G}_2$:

$$\mathbf{P}_k : \mathbb{R}^m \mapsto \text{Ker } \mathbf{G}_k, \quad \mathbf{P}_k := \mathbf{I} - \mathbf{G}_k^\top (\mathbf{G}_k \mathbf{G}_k^\top)^{-1} \mathbf{G}_k, \quad k = 1, 2$$

Property: $\text{Ker } \mathbf{P}_k = \text{Im } \mathbf{G}_k^\top \iff \mathbf{P}_k \mathbf{G}_k^\top = \mathbf{0}$

• \mathbf{P}_1 splits the saddle-point structure: $\mathbf{P}_1 \mathbf{F} \boldsymbol{\lambda} + \mathbf{P}_1 \mathbf{G}_1^\top \boldsymbol{\alpha} = \mathbf{P}_1 \mathbf{d}$

$$\mathbf{P}_1 \mathbf{F} \boldsymbol{\lambda} = \mathbf{P}_1 \mathbf{d}, \quad \mathbf{G}_2 \boldsymbol{\lambda} = \mathbf{e}, \quad \boldsymbol{\alpha} := (\mathbf{G}_1 \mathbf{G}_1^\top)^{-1} (\mathbf{G}_1 \mathbf{d} - \mathbf{G}_1 \mathbf{F} \boldsymbol{\lambda})$$

• \mathbf{P}_2 decomposes $\boldsymbol{\lambda} = \boldsymbol{\lambda}_{\text{Im}} + \boldsymbol{\lambda}_{\text{Ker}}$, $\boldsymbol{\lambda}_{\text{Im}} \in \text{Im } \mathbf{G}_2^\top$, $\boldsymbol{\lambda}_{\text{Ker}} \in \text{Ker } \mathbf{G}_2$

At first: $\mathbf{G}_2 \boldsymbol{\lambda} = \mathbf{G}_2 \boldsymbol{\lambda}_{\text{Im}} = \mathbf{e} \implies \boldsymbol{\lambda}_{\text{Im}} := \mathbf{G}_2^\top (\mathbf{G}_2 \mathbf{G}_2^\top)^{-1} \mathbf{e}$

At second: $\mathbf{P}_1 \mathbf{F} \boldsymbol{\lambda}_{\text{Ker}} = \mathbf{P}_1 (\mathbf{d} - \mathbf{F} \boldsymbol{\lambda}_{\text{Im}})$ on $\text{Ker } \mathbf{G}_2$

Theorem 4 Let \mathcal{A} be invertible. The linear operator $\mathbf{P}_1\mathbf{F}: Ker \mathbf{G}_2 \mapsto Ker \mathbf{G}_1$ is invertible.

Proof.

As both null-spaces $Ker \mathbf{G}_1$ and $Ker \mathbf{G}_2$ have the same dimension $m - p$, it is enough to prove that $\mathbf{P}_1\mathbf{F}$ is injective.

Let $\boldsymbol{\mu} \in Ker \mathbf{G}_2$ be such that $\mathbf{P}_1\mathbf{F}\boldsymbol{\mu} = \mathbf{0}$. Then $\mathbf{F}\boldsymbol{\mu} \in Ker \mathbf{P}_1 = Im \mathbf{G}_1^\top$ and, therefore, there is $\boldsymbol{\beta} \in \mathbb{R}^p$ so that

$$\mathbf{F}\boldsymbol{\mu} = \mathbf{G}_1^\top\boldsymbol{\beta} \quad \text{and} \quad \mathbf{G}_2\boldsymbol{\mu} = \mathbf{0}.$$

We obtain

$$\begin{pmatrix} \mathbf{F} & \mathbf{G}_1^\top \\ \mathbf{G}_2 & \mathbf{0} \end{pmatrix} \begin{pmatrix} \boldsymbol{\mu} \\ -\boldsymbol{\beta} \end{pmatrix} = \begin{pmatrix} \mathbf{0} \\ \mathbf{0} \end{pmatrix},$$

where the matrix is the (negative) Schur complement $-\mathcal{S}$ that is invertible iff \mathcal{A} is invertible. Therefore $\boldsymbol{\mu} = \mathbf{0}$.

Algorithm PSCM

Step 1.a: Assemble $\mathbf{G}_1 := -\mathbf{N}^\top \mathbf{B}_2^\top$, $\mathbf{G}_2 := -\mathbf{M}^\top \mathbf{B}_1^\top$.

Step 1.b: Assemble $\mathbf{d} := \mathbf{B}_2 \mathbf{A}^\dagger \mathbf{f} - \mathbf{g}$, $\mathbf{e} := -\mathbf{M}^\top \mathbf{f}$.

Step 1.c: Assemble $\mathbf{H}_1 := (\mathbf{G}_1 \mathbf{G}_1^\top)^{-1}$, $\mathbf{H}_2 := (\mathbf{G}_2 \mathbf{G}_2^\top)^{-1}$.

Step 1.d: Assemble $\lambda_{Im} := \mathbf{G}_2^\top \mathbf{H}_2 \mathbf{e}$, $\tilde{\mathbf{d}} := \mathbf{P}_1(\mathbf{d} - \mathbf{F} \lambda_{Im})$.

Step 1.e: Solve $\mathbf{P}_1 \mathbf{F} \lambda_{Ker} = \tilde{\mathbf{d}}$ on $Ker \mathbf{G}_2$.

Step 1.f: Assemble $\lambda := \lambda_{Im} + \lambda_{Ker}$.

Step 2: Assemble $\alpha := \mathbf{H}_1 \mathbf{G}_1(\mathbf{d} - \mathbf{F} \lambda)$.

Step 3: Assemble $\mathbf{u} := \mathbf{A}^\dagger(\mathbf{f} - \mathbf{B}_1^\top \lambda) + \mathbf{N} \alpha$.

- an iterative **projected** Krylov subspace method for non-symmetric operators can be used in Step 1.e

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Find $\lambda \in \mathbb{R}^m$ so that $\mathbf{F}\lambda = \mathbf{d}$, where $\mathbf{d} \in \mathbb{R}^m$.

Algorithm BiCGSTAB $[\epsilon, \lambda^0, \mathbf{F}, \mathbf{d}] \rightarrow \lambda$

Initialize: $\mathbf{r}^0 := \mathbf{d} - \mathbf{F}\lambda^0$, $\mathbf{p}^0 := \mathbf{r}^0$, $\tilde{\mathbf{r}}^0$ arbitrary, $k := 0$

While $\|\mathbf{r}^k\| > \epsilon$

1° $\tilde{\mathbf{p}}^k := \mathbf{F}\mathbf{p}^k$

2° $\alpha_k := (\mathbf{r}^k)^\top \tilde{\mathbf{r}}^0 / (\tilde{\mathbf{p}}^k)^\top \tilde{\mathbf{r}}^0$

3° $\mathbf{s}^k := \mathbf{r}^k - \alpha_k \tilde{\mathbf{p}}^k$

4° $\tilde{\mathbf{s}}^k := \mathbf{F}\mathbf{s}^k$

5° $\omega_k := (\tilde{\mathbf{s}}^k)^\top \mathbf{s}^k / (\tilde{\mathbf{s}}^k)^\top \tilde{\mathbf{s}}^k$

6° $\lambda^{k+1} := \lambda^k + \alpha_k \mathbf{p}^k + \omega_k \mathbf{s}^k$

7° $\mathbf{r}^{k+1} := \mathbf{s}^k - \omega_k \tilde{\mathbf{s}}^k$

8° $\beta_{k+1} := (\alpha_k / \omega_k) (\mathbf{r}^{k+1})^\top \tilde{\mathbf{r}}^0 / (\mathbf{r}^k)^\top \tilde{\mathbf{r}}^0$

9° $\mathbf{p}^{k+1} := \mathbf{r}^{k+1} + \beta_{k+1} (\mathbf{p}^k - \omega_k \tilde{\mathbf{p}}^k)$

10° $k := k + 1$

end

(Van der Vorst, 1992)

Find $\lambda \in Ker \mathbf{G}_2$ so that $\mathbf{P}_1 \mathbf{F} \lambda = \tilde{\mathbf{d}}$, where $\tilde{\mathbf{d}} \in Ker \mathbf{G}_1$.

Algorithm ProjBiCGSTAB $[\epsilon, \lambda^0, \mathbf{F}, \mathbf{P}_1, \mathbf{P}_2, \tilde{\mathbf{d}}] \rightarrow \lambda$

Initialize: $\lambda^0 \in Ker \mathbf{G}_2$, $\mathbf{r}^0 := \tilde{\mathbf{d}} - \mathbf{P}_1 \mathbf{F} \lambda^0$, $\mathbf{p}^0 := \mathbf{r}^0$, $\tilde{\mathbf{r}}^0$ arbitrary, $k := 0$

While $\|\mathbf{r}^k\| > \epsilon$

1° $\tilde{\mathbf{p}}^k := \mathbf{P}_1 \mathbf{F} \mathbf{p}^k$

2° $\alpha_k := (\mathbf{r}^k)^\top \tilde{\mathbf{r}}^0 / (\tilde{\mathbf{p}}^k)^\top \tilde{\mathbf{r}}^0$

3° $\mathbf{s}^k := \mathbf{r}^k - \alpha_k \tilde{\mathbf{p}}^k$

4° $\tilde{\mathbf{s}}^k := \mathbf{P}_1 \mathbf{F} \mathbf{s}^k$

5° $\omega_k := (\tilde{\mathbf{s}}^k)^\top \mathbf{s}^k / (\tilde{\mathbf{s}}^k)^\top \tilde{\mathbf{s}}^k$

6° $\lambda^{k+1} := \lambda^k + \alpha_k \mathbf{P}_2 \mathbf{p}^k + \omega_k \mathbf{P}_2 \mathbf{s}^k$

7° $\mathbf{r}^{k+1} := \mathbf{s}^k - \omega_k \tilde{\mathbf{s}}^k$

8° $\beta_{k+1} := (\alpha_k / \omega_k) (\mathbf{r}^{k+1})^\top \tilde{\mathbf{r}}^0 / (\mathbf{r}^k)^\top \tilde{\mathbf{r}}^0$

9° $\mathbf{p}^{k+1} := \mathbf{r}^{k+1} + \beta_{k+1} (\mathbf{p}^k - \omega_k \tilde{\mathbf{p}}^k)$

10° $k := k + 1$

end

Formally solve $\mathbf{P}_2\mathbf{P}_1\mathbf{F}\boldsymbol{\lambda} = \mathbf{P}_2\tilde{\mathbf{d}}$, with $\boldsymbol{\lambda}^0 \in \text{Ker } \mathbf{G}_2$.

Algorithm ProjBiCGSTAB $[\epsilon, \boldsymbol{\lambda}^0, \mathbf{F}, \mathbf{P}_1, \mathbf{P}_2, \tilde{\mathbf{d}}] \rightarrow \boldsymbol{\lambda}$

Initialize: $\boldsymbol{\lambda}^0 \in \text{Ker } \mathbf{G}_2$, $\mathbf{r}^0 := \mathbf{P}_2\tilde{\mathbf{d}} - \mathbf{P}_2\mathbf{P}_1\mathbf{F}\boldsymbol{\lambda}^0$, $\mathbf{p}^0 := \mathbf{r}^0$, $\tilde{\mathbf{r}}^0$, $k := 0$

While $\|\mathbf{r}^k\| > \epsilon$

1° $\tilde{\mathbf{p}}^k := \mathbf{P}_2\mathbf{P}_1\mathbf{F}\mathbf{p}^k$

2° $\alpha_k := (\mathbf{r}^k)^\top \tilde{\mathbf{r}}^0 / (\tilde{\mathbf{p}}^k)^\top \tilde{\mathbf{r}}^0$

3° $\mathbf{s}^k := \mathbf{r}^k - \alpha_k \tilde{\mathbf{p}}^k$

4° $\tilde{\mathbf{s}}^k := \mathbf{P}_2\mathbf{P}_1\mathbf{F}\mathbf{s}^k$

5° $\omega_k := (\tilde{\mathbf{s}}^k)^\top \mathbf{s}^k / (\tilde{\mathbf{s}}^k)^\top \tilde{\mathbf{s}}^k$

6° $\boldsymbol{\lambda}^{k+1} := \boldsymbol{\lambda}^k + \alpha_k \mathbf{p}^k + \omega_k \mathbf{s}^k$

7° $\mathbf{r}^{k+1} := \mathbf{s}^k - \omega_k \tilde{\mathbf{s}}^k$

8° $\beta_{k+1} := (\alpha_k / \omega_k) (\mathbf{r}^{k+1})^\top \tilde{\mathbf{r}}^0 / (\mathbf{r}^k)^\top \tilde{\mathbf{r}}^0$

9° $\mathbf{p}^{k+1} := \mathbf{r}^{k+1} + \beta_{k+1} (\mathbf{p}^k - \omega_k \tilde{\mathbf{p}}^k)$

10° $k := k + 1$

end

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Numerical experiments

Consider a family of nested partitions of the fictitious domain Ω with stepsizes:

$$h_j, \quad 0 \leq j \leq J$$

- the first iterate is determined by the result from the nearest lower level
- the terminating tolerance ϵ on each level is $\epsilon := \nu h_j^p$

Algorithm: Hierarchical Multigrid Scheme

Initialize: Let $\lambda_{Ker}^{0,(0)} \in Ker \mathbf{G}_2^{(0)}$ be given.

ProjBiCGSTAB[$\nu h_0^p, \lambda_{Ker}^{0,(0)}, \mathbf{F}^{(0)}, \mathbf{P}_1^{(0)}, \mathbf{P}_2^{(0)}, \tilde{\mathbf{d}}^{(0)}$] $\rightarrow \lambda_{Ker}^{(0)}$.

For $j = 1, \dots, J$,

1° prolongate $\lambda_{Ker}^{(j-1)} \rightarrow \tilde{\lambda}_{Ker}^{0,(j)}$

2° project $\tilde{\lambda}_{Ker}^{0,(j)} \rightarrow \lambda_{Ker}^{0,(j)} := \mathbf{P}_2^{(j)} \tilde{\lambda}_{Ker}^{0,(j)}$

3° ProjBiCGSTAB[$\nu h_j^p, \lambda_{Ker}^{0,(j)}, \mathbf{F}^{(j)}, \mathbf{P}_1^{(j)}, \mathbf{P}_2^{(j)}, \tilde{\mathbf{d}}^{(j)}$] $\rightarrow \lambda_{Ker}^{(j)}$

end

Return: $\lambda_{Ker} := \lambda_{Ker}^{(J)}$.

Motivation

\mathbf{u}^* ... exact solution of PDE problem

\mathbf{u} ... FEM approximation with respect to h with the convergence rate p

$$\|\mathbf{u}^* - \mathbf{u}\| \leq Ch^p, \quad \mathbf{A}\mathbf{u} = \mathbf{f}$$

\mathbf{u}^k ... the k -th iteration

$$\mathbf{u}^k \longrightarrow \mathbf{u}, \quad \mathbf{A}\mathbf{u}^k = \mathbf{f} + \mathbf{r}^k$$

When should be iterations terminated? $\|\mathbf{r}^k\| \leq \epsilon$, $\epsilon = ???$

$$\begin{aligned} \|\mathbf{u}^* - \mathbf{u}^k\| &\leq \|\mathbf{u}^* - \mathbf{u}\| + \|\mathbf{u} - \mathbf{u}^k\| \\ &\leq Ch^p + \|\mathbf{A}^{-1}\mathbf{r}^k\| \\ &\leq Ch^p + \|\mathbf{A}^{-1}\| \cdot \epsilon \\ &\leq (C + \|\mathbf{A}^{-1}\|\nu)h^p \quad \text{if } \epsilon := \nu h^p \end{aligned}$$

Control parameter ν may be chosen experimentally; $\nu \approx KC/\|\mathbf{A}^{-1}\|$.

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Numerical experiments

Circulant matrices and Fourier transform

$$\mathbf{A} = \begin{pmatrix} a_1 & a_n & \dots & a_2 \\ a_2 & a_1 & \dots & a_3 \\ a_3 & a_2 & \dots & a_4 \\ \vdots & \vdots & \ddots & \vdots \\ a_n & a_{n-1} & \dots & a_1 \end{pmatrix} = (\mathbf{a}, \mathbf{T}\mathbf{a}, \mathbf{T}^2\mathbf{a}, \dots, \mathbf{T}^{n-1}\mathbf{a})$$

$$\widehat{\mathcal{T}_k f}(\omega) = \int_R f(x - k) e^{-ix\omega} dx = e^{-ik\omega} \widehat{f}(\omega)$$

$$\mathbf{X}\mathbf{A} = (\mathbf{D}\mathbf{x}_0, \mathbf{D}\mathbf{x}_1, \mathbf{D}\mathbf{x}_2, \dots, \mathbf{D}\mathbf{x}_{n-1}) = \mathbf{D}\mathbf{X}$$

Lamma: Let \mathbf{A} be circulant. Then

$$\mathbf{A} = \mathbf{X}^{-1}\mathbf{D}\mathbf{X},$$

where \mathbf{X} is the DFT matrix and $\mathbf{D} = \text{diag}(\widehat{\mathbf{a}})$, $\widehat{\mathbf{a}} = \mathbf{X}\mathbf{a}$, $\mathbf{a} = \mathbf{A}(:, 1)$.

Multiplying procedure: $\mathbf{A}^\dagger \mathbf{v} := \mathbf{X}^{-1} (\mathbf{D}^\dagger (\mathbf{X}\mathbf{v})) \quad \dots \quad \text{Moore-Penrose}$

$$\left. \begin{array}{l} 0^\circ \quad \mathbf{d} := \text{fft}(\mathbf{a}) \\ 1^\circ \quad \mathbf{v} := \text{fft}(\mathbf{v}) \\ 2^\circ \quad \mathbf{v} := \mathbf{v} * \mathbf{d}^{-1} \\ 3^\circ \quad \mathbf{A}^\dagger \mathbf{v} := \text{ifft}(\mathbf{v}) \end{array} \right\} \mathcal{O}(2n \log_2 n)$$

Multiplying procedures: $\mathbf{N}\boldsymbol{\alpha}$, $\mathbf{N}^\top \mathbf{v}$ (and $\mathbf{M}\boldsymbol{\alpha}$, $\mathbf{M}^\top \mathbf{v}$)

As $\mathbf{A}\mathbf{N} = \mathbf{0}$, the matrix \mathbf{N} may be formed by eigenvectors corresponding to zero eigenvalues.

$$\mathbf{I} - \mathbf{D}\mathbf{D}^\dagger = \text{diag}(1, 1, 1, 0, \dots, 0) \implies \mathbf{X}^{-1} = (\mathbf{N}, \mathbf{Y}), \quad \mathbf{X} = \begin{pmatrix} \mathbf{N}^\top \\ \mathbf{Y} \end{pmatrix}$$

Therefore we can define the operation: $\text{ind}(\boldsymbol{\alpha}) = \begin{pmatrix} \boldsymbol{\alpha} \\ \mathbf{0} \end{pmatrix} \in \mathbb{R}^n$

$$\left. \begin{array}{ll} 1^\circ \quad \mathbf{v}_\alpha := \text{ind}(\boldsymbol{\alpha}) & 1^\circ \quad \mathbf{v} := \text{ifft}(\mathbf{v}) \\ 2^\circ \quad \mathbf{N}\boldsymbol{\alpha} := \text{ifft}(\mathbf{v}_\alpha) & 2^\circ \quad \mathbf{N}^\top \mathbf{v} := \text{ind}^{-1}(\mathbf{v}) \end{array} \right\} \mathcal{O}(n \log_2 n)$$

Kronecker product of matrices: $\mathbf{A}_x \in \mathbb{R}^{n_x \times n_x}$, $\mathbf{A}_y \in \mathbb{R}^{n_y \times n_y}$

$$\mathbf{A}_x \otimes \mathbf{A}_y = \begin{pmatrix} a_{11}^y \mathbf{A}_x & \dots & a_{1n_y}^y \mathbf{A}_x \\ \vdots & \ddots & \vdots \\ a_{n_y 1}^y \mathbf{A}_x & \dots & a_{n_y n_y}^y \mathbf{A}_x \end{pmatrix}$$

Lemma 1:

$$\begin{aligned} (\mathbf{A}_x \otimes \mathbf{A}_y)(\mathbf{B}_x \otimes \mathbf{B}_y) &= \mathbf{A}_x \mathbf{B}_x \otimes \mathbf{A}_y \mathbf{B}_y \\ (\mathbf{A}_x \otimes \mathbf{A}_y)^\dagger &= \mathbf{A}_x^\dagger \otimes \mathbf{A}_y^\dagger \\ \mathbf{N} &= \mathbf{N}_x \otimes \mathbf{N}_y \end{aligned}$$

Lemma 2: $(\mathbf{A}_x \otimes \mathbf{A}_y)\mathbf{v} = \text{vec}(\mathbf{A}_x \mathbf{V} \mathbf{A}_y^\top)$, where $\mathbf{V} = \text{vec}^{-1}(\mathbf{v})$.

$$\mathbf{V} = (\mathbf{v}_1, \dots, \mathbf{v}_{n_y}) \in \mathbb{R}^{n_x \times n_y} \iff \text{vec}(\mathbf{V}) = \begin{pmatrix} \mathbf{v}_1 \\ \vdots \\ \mathbf{v}_{n_y} \end{pmatrix} \in \mathbb{R}^{n_x n_y}$$

Kronecker product and circulant matrices: Let $\mathbf{A}_x, \mathbf{A}_y$ be circulant then:

$$\begin{aligned}\mathbf{A} &= \mathbf{A}_x \otimes \mathbf{I}_y + \mathbf{I}_x \otimes \mathbf{A}_y \\ &= \mathbf{X}_x^{-1} \mathbf{D}_x \mathbf{X}_x \otimes \mathbf{X}_y^{-1} \mathbf{X}_y + \mathbf{X}_x^{-1} \mathbf{X}_x \otimes \mathbf{X}_y^{-1} \mathbf{D}_y \mathbf{X}_y \\ &= (\mathbf{X}_x^{-1} \otimes \mathbf{X}_y^{-1}) (\mathbf{D}_x \otimes \mathbf{I}_y + \mathbf{I}_x \otimes \mathbf{D}_y) (\mathbf{X}_x \otimes \mathbf{X}_y) \\ &= \mathbf{X}^{-1} \mathbf{D} \mathbf{X}\end{aligned}$$

with

$$\mathbf{X} = \mathbf{X}_x \otimes \mathbf{X}_y \quad (\text{DFT matrix in 2D})$$

$$\mathbf{D} = \mathbf{D}_x \otimes \mathbf{I}_y + \mathbf{I}_x \otimes \mathbf{D}_y \quad (\text{diagonal matrix})$$

where $\mathbf{X}_x, \mathbf{X}_y$ are the DFT matrices, $\mathbf{D}_x = \text{diag}(\mathbf{X}_x \mathbf{a}_x)$, $\mathbf{D}_y = \text{diag}(\mathbf{X}_y \mathbf{a}_y)$ and $\mathbf{a}_x = \mathbf{A}_x(:, 1)$, $\mathbf{a}_y = \mathbf{A}_y(:, 1)$, respectively.

Multiplying procedure: $\mathbf{A}^\dagger \mathbf{v} := \mathbf{X}^{-1} (\mathbf{D}^\dagger (\mathbf{X} \mathbf{v}))$

$$0^\circ \quad \mathbf{d}_x := \text{fft}(\mathbf{a}_x), \quad \mathbf{d}_y := \text{fft}(\mathbf{a}_y)$$

$$\mathbf{V} := \text{vec}^{-1}(\mathbf{v})$$

$$1^\circ \quad \mathbf{V} := \text{fft}(\mathbf{V})$$

$$2^\circ \quad \mathbf{V} := \text{fft}(\mathbf{V}^\top)^\top$$

$$3^\circ \quad \mathbf{V} := \text{vec}^{-1}(\mathbf{D}^\dagger \text{vec}(\mathbf{V}))$$

$$4^\circ \quad \mathbf{V} := \text{ifft}(\mathbf{V})$$

$$5^\circ \quad \mathbf{V} := \text{ifft}(\mathbf{V}^\top)^\top$$

$$\mathbf{A}^\dagger \mathbf{v} := \text{vec}(\mathbf{V})$$

Number of arithmetic operations :

$$\mathcal{O}(2n(\log_2 n_x + \log_2 n_y) + n) \approx \mathcal{O}(n \log_2 n), \quad n = n_x n_y$$

Multiplying procedures: $\mathbf{N}\boldsymbol{\alpha}, \mathbf{N}^\top \mathbf{v}, \mathbf{M}\boldsymbol{\alpha}, \mathbf{M}^\top \mathbf{v} \quad \dots \quad \text{analogous}$

OUTLINE

Motivation: Fictitious domain method

Algorithm PSCM: Schur complement method + Null-space method

Inner solver: Projected BiCGSTAB

Preconditioning: Hierarchical multigrid

Singular matrices: Poisson-like solver based on circulants

Numerical experiments

CONCLUSIONS

- The method for solving non-symmetric saddle-point systems with singular diagonal blocks was presented. It combines the Schur complement reduction with the null-space method.
- It can be understood as a generalization of the algebraic description of FETI DDM for non-symmetric and possibly indefinite cases.
- In connection with FDM, it presents the highly efficient solver for solving separable PDE problems. The fast implementation based on the Poisson-like solver is "matrix free" as the stiffness matrix is not needed to be formed explicitly.

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