

## A Step Size Rule for Unconstrained Optimization

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**Abstract.** We propose an Armijo-type step size rule for solving differentiable unconstrained optimization problems, based on a quadratic approximation. The rule is proved to be finite and to perform the exact line search in one iteration in case of a strictly convex quadratic function. We give convergence results for the steepest descent and for the FR, PR and DFP methods using this rule. Computational evidence shows that the rule performs best when implemented into the DFP or BFGS method.

**Key words.** Unconstrained optimization, step size rule, convergence

### 1. Introduction

Gradient methods for solving an unconstrained optimization problem

$$\begin{aligned} & \text{minimize } f(x) \\ & \text{subject to } x \in \mathbb{R}^n, \end{aligned}$$

as e.g. the steepest descent (SD) method or the methods by Fletcher-Reeves (FR), Polak-Ribière (PR), Davidon-Fletcher-Powell (DFP) or Broyden-Fletcher-Goldfarb-Shanno (BFGS), described in standard textbooks like Luenberger [1], Fletcher [2] or Polak [3], construct a sequence of iterations  $\{x_1\}$  according to this general scheme (which we call the "main algorithm" to distinguish it from its concrete specifications; we denote  $g_1 = \nabla f(x_1)$ , the gradient of  $f$  at  $x_1$ ):

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Main algorithm.

Step\_0. Select an  $x_0 \in R^n$  and set  $i:=0$ .

Step\_1. If  $g_i = 0$ , terminate:  $x_i$  is a stationary point of  $f$ .

Step\_2. Otherwise find a search direction  $d_i$  such that  $d_i^T g_i < 0$ .

Step\_3. Find a real number  $\alpha_i$  satisfying

$$f(x_i + \alpha_i d_i) = \min \{f(x_i + \alpha d_i); \alpha \geq 0\}.$$

Step\_4. Set  $x_{i+1} := x_i + \alpha_i d_i$ ,  $i := i + 1$  and go to Step 1.

The methods listed above differ from each other only in the choice of the search direction  $d_i$  in Step 2 (to be described later in Section 3). All of them share the common descent property of the main algorithm, namely that  $f(x_{i+1}) < f(x_i)$  for each  $i$ . The procedure of minimizing the function of one real variable  $f(x_i + \alpha d_i)$  over the nonnegative half-ray is called the line search, and the value  $\alpha_i$  the step size. In the form stated the main algorithm remains, however, only of conceptual value since generally the exact minimum in Step 3 cannot be found in a finite number of iterations; therefore in practical computations the exact line search in Step 3 must be substituted by some finite procedure (a so-called inexact line search) yielding satisfactory approximate results. Among many options discussed in [1], [2], [3], we shall focus our attention here on the step size rule proposed by Armijo [4], using a preselected parameter  $\beta \in (0,1)$ :

Step\_3A. Set  $\alpha_i := \beta^k$ , where  $k$  is the minimum positive integer satisfying

$$f(x_i + \beta^k d_i) - f(x_i) \leq 0.5 \beta^k d_i^T g_i .$$

In view of the choice of  $d_i$  made in Step 2, it is easy to see that: (a) such a  $k$  exists, and (b) the sequence generated by the main algorithm using this rule satisfies  $f(x_{i+1}) < f(x_i)$  for each  $i$ , hence

the descent property is not violated.

In this paper we shall propose another rule, which is in certain sense similar to that of Armijo's, but based on quadratic considerations. In section 2 we formulate the rule and prove that it is finite and maintains the descent property; moreover, if  $f$  is a strictly convex quadratic function, then the line search is exact and the minimizing value is always found in one iteration. In section 3 we state a general convergence result (Theorem 3.1) which is then applied to establish convergence properties of the SD, FR, PR and DFP methods using this rule. Computational experience shows that the rule performs well, especially when implemented into the DFP or BFGS method.

## 2. The rule

We shall first give an informal description of the basic idea. Assume at the  $i$ -th iteration of the main algorithm we have an approximation  $\beta_j > 0$  ( $j = 0, 1, \dots$ ) of the exact step size. Assuming  $f$  to be at least twice continuously differentiable, let us replace the function  $f(x_1 + \beta d_1)$  of one real variable  $\beta$  by its truncated Taylor series

$$f(x_1 + \beta d_1) \doteq f(x_1) + \beta d_1^T g_1 + 0.5 \beta^2 d_1^T H_1 d_1 \quad (1)$$

where  $H_1 = (\partial^2 f(x_1) / \partial x_j \partial x_k)_{j,k=1}^n$  is the Hessian matrix at  $x_1$ . It is obvious that the right-hand side quadratic function achieves its minimum (if it exists) at the point

$$\beta = -(d_1^T g_1) / (d_1^T H_1 d_1). \quad (2)$$

To avoid the necessity of evaluating  $H_1$ , let us express the denominator in (2) from the formula (1) applied to the previously computed approximation  $\beta_j$ :

$$d_1^T H_1 d_1 \doteq (2/\beta_j^2)(f(x_1 + \beta_j d_1) - f(x_1) - \beta_j d_1^T g_1). \quad (3)$$

Hence, using (2) and (3), we may construct the next approximation by

$$\beta_{j+1} = -0.5 \beta_j^2 d_1^T g_1 / (f(x_1 + \beta_j d_1) - f(x_1) - \beta_j d_1^T g_1). \quad (4)$$

Now observe that this equation implies

$$f(x_1 + \beta_j d_1) - f(x_1) = (\beta_j - 0.5 \beta_j^2 / \beta_{j+1}) d_1^T g_1. \quad (5)$$

Thus, if

$$\beta_j - 0.5 \beta_j^2 / \beta_{j+1} > 0, \quad (6)$$

then  $f(x_1 + \beta_j d_1) - f(x_1) < 0$ , so that we may terminate the procedure and set  $\alpha_1 := \beta_j$ , which will enforce the descent property

$$f(x_{1+1}) < f(x_1) \quad (7)$$

to hold. Since  $\beta_j > 0$ , we may simplify (6) to

$$\beta_j / \beta_{j+1} < 2.$$

In the opposite case

$$\beta_j / \beta_{j+1} \geq 2$$

we have  $\beta_{j+1} > 0$  and setting  $j := j+1$ , we may repeat the whole process.

Finally, if the denominator in (4) is zero, then we have

$$f(x_1 + \beta_j d_1) - f(x_1) = \beta_j d_1^T g_1 < 0, \quad (8)$$

so that we can set  $\alpha_1 := \beta_j$ . Summing up, the line search procedure runs as follows:

Step Size Rule 2.1 (to replace Step 3 of the main algorithm).

Step 3.1. Set  $\beta_0 := 1$  and  $j := 0$ .

Step 3.2. Compute  $\gamma_j = f(x_1 + \beta_j d_1) - f(x_1) - \beta_j d_1^T g_1$ .

Step 3.3. If  $\gamma_j = 0$ , set  $\alpha_1 := \beta_j$  and go to Step 4.

Step 3.4. Otherwise compute  $\beta_{j+1} = -0.5(\beta_j^2 / \gamma_j) d_1^T g_1$ .

Step 3.5. If  $\beta_j / \beta_{j+1} < 2$ , set  $\alpha_1 := \beta_j$  and go to Step 4.

Step 3.6. Otherwise (i.e. if  $\beta_j / \beta_{j+1} \geq 2$ ) set  $j := j+1$  and go to Step 3.2.

It follows from Step 3.6 that if the rule does not stop with  $\beta_j$ , then the newly constructed iteration satisfies

$$0 < \beta_{j+1} \leq 0.5 \beta_j. \quad (9)$$

We shall now prove that  $\omega_1$  is reached after a finite number of iterations:

Proposition 2.1. Let  $f \in C^1$ . Then the step size rule 2.1 is finite and the main algorithm using this rule generates a sequence of points satisfying

$$f(x_{i+1}) < f(x_i)$$

for each  $i$ .

Proof. Assume to the contrary that the rule never terminates for some  $i$ , so that it constructs a sequence of numbers  $\{\beta_j\}$  tending to zero in view of (9). Hence from (5) we have

$$(f(x_i + \beta_j d_i) - f(x_i)) / \beta_j = (1 - 0.5 \beta_j / \beta_{j+1}) d_i^T g_i$$

for each  $j$ . Since the left-hand side tends to  $d_i^T g_i$  as  $j$  approaches infinity, we obtain

$$\lim_{j \rightarrow \infty} \beta_j / \beta_{j+1} = 0.$$

But this is a contradiction since  $\beta_j / \beta_{j+1} \geq 2$  for each  $j$  due to Step 3.6. Hence the rule is finite. The fact that the objective strictly decreases has been already established in (7) and (8).

Proposition 2.2. If  $f$  is a strictly convex quadratic function, then the line search 3.1-3.6 is exact and  $\omega_1 = \beta_1$  for each  $i$ .

Proof. Let  $f(x) = 0.5x^T Cx + b^T x + a$ , where  $C$  is a symmetric positive definite matrix. Then for each  $\beta$  we have

$$f(x_i + \beta d_i) = f(x_i) + \beta d_i^T g_i + 0.5 \beta^2 d_i^T C d_i,$$

hence  $\gamma_j > 0$  since  $C$  is positive definite and  $d_1 \neq 0$  by Step 2, and for each  $j$  we have from (4)

$$\beta_{j+1} = -(d_1^T g_j) / (d_1^T C d_1),$$

the right-hand side constant being obviously the exact minimizer of  $f(x_1 + \beta d_1)$  over the nonnegative half-ray. Hence  $\beta_2 = \beta_1$ , so that  $\omega_1 := \beta_1$  is set in Step 3.5.

### 3. Convergence properties

In this section we shall first prove in Theorem 3.1 a convergence property for the main algorithm endowed with the step size rule 2.1 and then will use it to establish some convergence results for the SD, FR, PR and DFP methods under the rule 2.1.

For each  $i = 0, 1, \dots$ , denote by  $j_i$  the index  $j$  for which  $\omega_i := \beta_j$  is set in Step 3.3 or Step 3.5 of the rule 2.1. If  $j_i \geq 1$ , denote also  $\bar{\omega}_i = \beta_{j_i-1}$ . Then for each  $i$  with  $j_i \geq 1$  we have

$$f(x_i + \bar{\omega}_i d_i) - f(x_i) \geq 0 \quad (10)$$

$$f(x_i + \omega_i d_i) - f(x_i) < 0. \quad (11)$$

In fact, (11) follows from Proposition 2.1. If  $j := j_i \geq 1$ , then the rule did not terminate at  $j-1$ , which means that it must have been

$$\beta_{j-1} / \beta_j \geq 2$$

in Step 3.6, implying

$$f(x_i + \bar{\omega}_i d_i) - f(x_i) = \beta_{j-1} (1 - 0.5 \beta_{j-1} / \beta_j) d_i^T g_i \geq 0$$

due to (5), which proves (10).

Theorem 3.1.1. Let  $f \in C^1$  and let the sequence generated by the main algorithm using the step size rule 2.1 have the property

$$x_{i+1} - x_i \rightarrow 0. \quad (12)$$

Then each accumulation point  $(x_*, d_*)$  of the sequence  $\{(x_i, d_i)\}$  satisfies

$$d_*^T \nabla f(x_*) = 0. \quad (13)$$

Proof. Let  $x_i \xrightarrow{K} x_*$ ,  $d_i \xrightarrow{K} d_*$  along some subsequence  $K = \{0, 1, 2, \dots\}$  which may be chosen so that  $\{\alpha_i\}_{i \in K}$  is convergent (since  $0 \leq \alpha_i \leq \max\{\rho, \beta\}$  for each  $i$ ). Let  $\alpha_i \xrightarrow{K} \alpha$ . We shall distinguish two cases:

(a) Let  $\alpha > 0$ . Then from  $\alpha_i d_i = x_{i+1} - x_i \rightarrow 0$  we have  $d_* = 0$ , hence (13) holds.

(b) Let  $\alpha = 0$ . Since  $d_i \xrightarrow{K} d_*$  and  $g_i \xrightarrow{K} \nabla f(x_*)$ , we have  $d_i^T g_i \xrightarrow{K} d_*^T \nabla f(x_*)$ . Assume to the contrary that  $d_*^T \nabla f(x_*) \neq 0$ . Since  $\alpha = 0$ , it must be  $j_i \geq 1$  from some  $i \in K$  on and from Step 3.4 we obtain

$$\alpha_i = -(\bar{\alpha}_i^2 / \bar{\gamma}_i) d_i^T g_i,$$

where  $\bar{\gamma}_i = \gamma_{j_i-1}$ . Here  $\alpha_i \xrightarrow{K} 0$ , the sequence  $\{\bar{\gamma}_i\}_K$  is bounded and  $d_i^T g_i \xrightarrow{K} d_*^T \nabla f(x_*) \neq 0$ , which gives  $\bar{\alpha}_i \xrightarrow{K} 0$ . Now from (10), (11) we have

$$\begin{aligned} f(x_i + \bar{\alpha}_i d_i) - f(x_i) &\geq 0 \\ f(x_i + \alpha_i d_i) - f(x_i) &< 0 \end{aligned}$$

for sufficiently large  $i \in K$ , hence by the continuity argument there is an  $\alpha'_i \in [\alpha_i, \bar{\alpha}_i]$  such that

$$f(x_i + \alpha'_i d_i) - f(x_i) = 0$$

holds and by the mean-value theorem we have

$$d_i^T \nabla f(x_i + \xi_i d_i) = 0$$

for some  $\xi_i \in [0, \alpha'_i] \subset [0, \bar{\alpha}_i]$ . Now taking the limit  $i \xrightarrow{K} \infty$ , from

$\overline{\omega}_1 \xrightarrow{K} 0$  we have  $\sum_1 \xrightarrow{K} 0$  and hence  

$$d_*^T \nabla f(x_*) = 0.$$

This contradicts the assumption  $d_*^T \nabla f(x_*) \neq 0$ . Hence (13) holds.

Comment. If the line search is exact (as assumed in the original Step 3 of the main algorithm), then taking the derivative of  $f(x_1 + \omega d_1)$  at the minimum point, we obtain

$$d_1^T g_{i+1} = 0 \quad (14)$$

for each  $i$  (cf. Polak [3]). If  $(x_1, d_1) \xrightarrow{K} (x_*, d_*)$  and (12) holds, then  $g_{i+1} \xrightarrow{K} \nabla f(x_*)$ , hence  $g_{i+1} - g_1 \xrightarrow{K} 0$  which gives  $d_1^T g_1 = -d_1^T (g_{i+1} - g_1) \xrightarrow{K} 0 = d_*^T \nabla f(x_*)$ . Hence the condition (13) is a generalization of (14).

We shall now apply Theorem 3.1 to establish convergence properties of several known methods under the step size rule 2.1; they arise by specifying the Step 2 of the main algorithm as follows (cf. [1]):

a) The steepest descent method:

$$d_1 = -g_1 \quad (i = 0, 1, \dots)$$

b) Conjugate gradient methods:

$$\begin{aligned} d_{i+1} &= -g_{i+1} + \delta_1 d_1 & (i = 0, 1, \dots) & \quad (15) \\ d_0 &= -g_0 \end{aligned}$$

where

$$\delta_1 = \|g_{i+1}\|_2^2 / \|g_1\|_2^2 \quad (\text{FR method}) \quad (16)$$

$$\delta_1 = g_{i+1}^T (g_{i+1} - g_1) / \|g_1\|_2^2 \quad (\text{PR method}). \quad (17)$$

c) The DFP method:

$$d_1 = -S_1 g_1 \quad (i = 0, 1, \dots) \quad (18)$$

where  $S_0$  is the unit matrix and updating is done by

$$S_{i+1} = S_i + (p_1 p_1^T) / (p_1^T q_1) - (S_i q_1 q_1^T S_i) / (q_1^T S_i q_1)$$

where  $p_1 = x_{i+1} - x_i$ ,  $q_1 = g_{i+1} - g_i$ ,  $i = 0, 1, \dots$ ; then all the matrices  $S_i$  are symmetric positive definite [1, p. 195].



Notice that the proofs which follow use only the property (13), not the actual construction of the sequence by the rule 2.1. Hence they are valid under any step size rule maintaining the property stated in Theorem 3.1.

Proposition 3.1. Let  $f \in C^1$  and let the sequence  $\{x_i\}$  constructed by the steepest descent method using the rule 2.1 satisfy (12). Then each accumulation point of  $\{x_i\}$  is a stationary point of  $f$ .

Proof. If  $x_i \xrightarrow{K} x_*$  along some  $K \subset \{0, 1, \dots\}$ , then  $d_i^T g_i = -\|g_i\|_2^2 \xrightarrow{K} -\|\nabla f(x_*)\|_2^2 = 0$  due to Theorem 3.1 and the assertion follows.

Proposition 3.2. Let  $f \in C^1$  and let both the sequences  $\{x_i\}$ ,  $\{d_i\}$  constructed by the FR (or PR) method using the rule 2.1 be convergent. Then the limit point of  $\{x_i\}$  is a stationary point of  $f$ .

Proof. Let  $x_i \rightarrow x_*$ ,  $d_i \rightarrow d_*$ , so that  $g_i \rightarrow g_* = \nabla f(x_*)$ . Then premultiplying (15) by  $g_{i+1}$ , from (16), (17) we have

$$\|g_i\|_2^2 (d_{i+1}^T g_{i+1}) = -\|g_i\|_2^2 \|g_{i+1}\|_2^2 + \bar{\delta}_i (d_i^T g_i + d_i^T (g_{i+1} - g_i)) \quad (19)$$

where  $\bar{\delta}_i = \|g_{i+1}\|_2^2$  for the FR method and  $\bar{\delta}_i = g_{i+1}^T (g_{i+1} - g_i)$  for the PR method. Since  $d_i^T g_i \rightarrow 0$  according to Theorem 3.1, taking the limit in (19) we obtain  $\|g_*\|_2^4 = 0$ , hence  $g_* = 0$ .

Proposition 3.3. Let  $f \in C^1$  and let  $\{x_i\}$  be constructed by the DFP method using the rule 2.1. Assume that (12) holds and that  $S_i \rightarrow S$ , where  $S$  is a positive definite matrix. Then each accumulation point of  $\{x_i\}$  is a stationary point of  $f$ .

Proof. If  $x_i \xrightarrow{K} x_*$ , then  $g_i \xrightarrow{K} g_* = \nabla f(x_*)$ , hence  $d_i \xrightarrow{K} -Sg_*$

by (18) and  $g_1^T S_1 g_1 \xrightarrow{K} g_*^T S g_*$ , but also  $g_1^T S_1 g_1 = -d_1^T g_1 \xrightarrow{K} 0$  by Theorem 3.1, hence  $g_*^T S g_* = 0$  and the positive definiteness of  $S$  implies  $g_* = 0$ .

Example. We have tested the rule 2.1 against the Armijo rule (with  $\beta = 0.7$ ) on the example used by Polak in [3], section C.2, for comparing the performances of various methods. The example reads

$$\min \{ \exp(x_1^2 + 5x_2^2) + x_1^2 + 80x_2^2; x \in \mathbb{R}^2 \}$$

and clearly has a unique optimal solution  $x = (0,0)^T$ . We have used throughout Polak's starting point  $x_0 = (1.32, -0.07)^T$  and the stopping rule  $\|x_k - x_{k-1}\|_\infty < 10^{-3}$ . Armijo rule gives the following results (at each method we depict the number of iterations  $k$  and the coordinates of the last iteration  $x_k$ ):

method	k	$(x_k)_1$	$(x_k)_2$
SD	35	2.79333E-02	3.60985E-04
FR	12	1.72954E-03	-1.88560E-05
PR	11	1.11621E-03	1.04664E-04
DFP	10	1.82042E-04	5.48999E-06
BFGS	9	1.97058E-04	-3.23192E-05

Using alternatively the rule 2.1, we obtained these results:

method	k	$(x_k)_1$	$(x_k)_2$
SD	22	3.56810E-02	-6.01563E-03
FR	10	-1.20377E-03	-1.37141E-05
PR	5	1.03677E-02	9.10693E-03
DFP	7	-5.23471E-07	-3.40049E-07
BFGS	6	-9.85223E-07	-5.75616E-07

It can be seen that the rule 2.1 performed approximately equally with SD and FR, a little worse with PR, but essentially better with DFP and BFGS. This particular result corresponds to our overall (although limited) computational experience according to which the rule performs best when incorporated into the DFP or BFGS method.

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